

# PARTIAL AUTOMORPHISMS OF COMBINATORIAL STRUCTURES

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# Combinatorial structure - classical approach to symmetries

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# Combinatorial structure - classical approach to symmetries

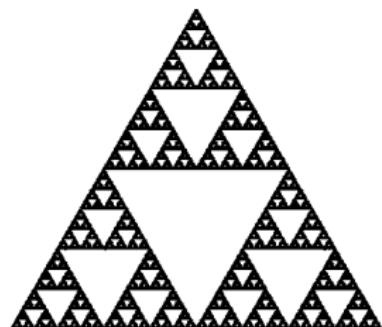
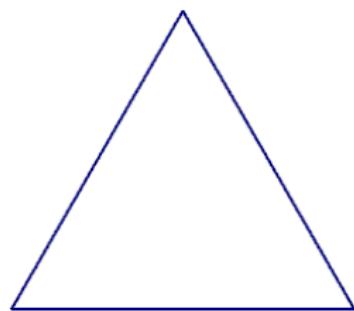
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Examples: graphs, directed graphs, hypergraphs, geometries, designs, ...

- ▶ symmetries - **automorphisms** of  $(V, \mathcal{F})$

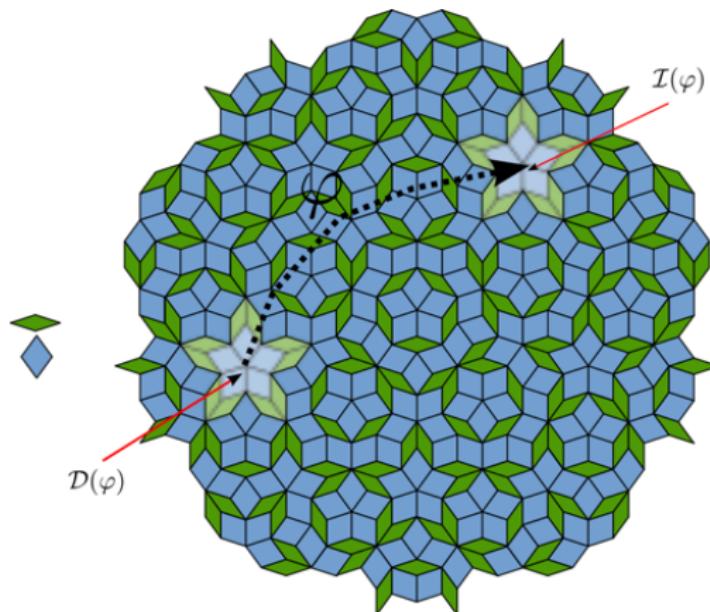
- ▶ Automorphisms form a group,  $\text{Aut}(\mathcal{C}) \leq \text{Sym}(V)$ .

# Partial Symmetries



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## Penrose tiling



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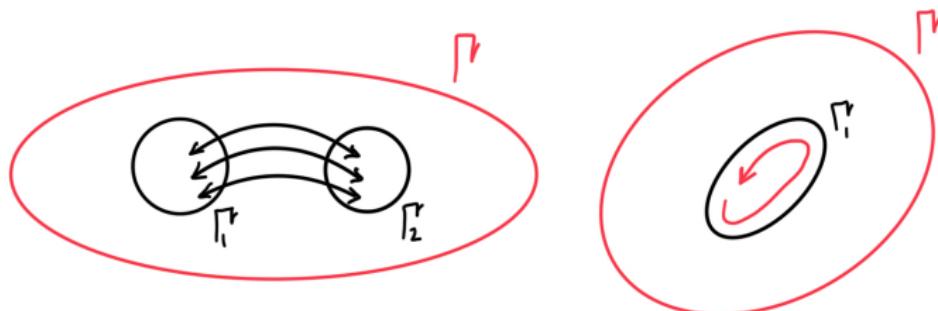
- ▶ (Erdős, Rényi, 1963)  
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- ▶ (Kim, Sudakov, Vu, 2002)  
Almost all regular graphs are asymmetric
- ▶ in particular,  $Aut(\Gamma)$  is trivial for all those graphs

# Partial Automorphisms

A **partial automorphism** of  $\Gamma = (V, \mathcal{E})$  is an isomorphism between two *induced* subgraphs.

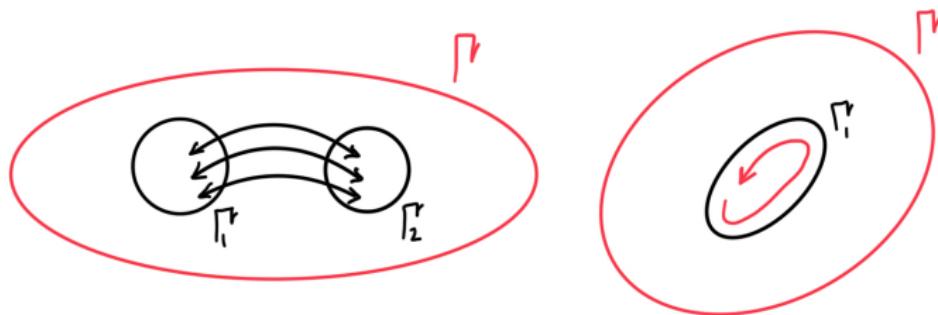
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A **partial automorphism** of a combinatorial structure  $\mathcal{C}$  is an isomorphism between two *induced* substructures.

# Inverse Monoid of Partial Automorphisms

The set of **all partial automorphisms**, denoted  $\text{PAut}(\Gamma)$  with the composition and partial inverse of partial maps forms an inverse monoid.

$$\text{PAut}(\Gamma) \leq \text{PSym}(V)$$

A **rank** of a partial automorphism is given by the size of its domain.

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- ▶ The inverse monoid  $\text{PAut}(\Gamma)$  is a complete algebraic description of  $\Gamma$
- ▶  $\text{PAut}(\Gamma)$  contains  $\text{Aut}(\Gamma)$  as its subgroup

# Wagner-Preston representation

While groups can be represented as **symmetries**:

## Theorem (Cayley)

*Every group can be embedded in the set of one to one transformations on a set.*

Inverse semigroups can be represented as **partial symmetries**:

## Theorem (Wagner-Preston)

*Every inverse semigroup can be embedded in the set of **partial** one to one transformations on a set.*

# Problem

Given a class of combinatorial structures and an inverse monoid  $M$ , is there a combinatorial structure  $\mathcal{C}$  from our class whose full partial automorphism monoid  $\text{PAut}(\mathcal{C})$  is isomorphic to  $M$ ?

Analogue of Frucht's theorem for groups.

# Classification Problem

Given a class of combinatorial structures, classify finite inverse monoids  $M$  for which there exists a structure  $\mathcal{C}$  from the considered class whose full partial automorphism monoid  $\text{PAut}(\mathcal{C})$  is isomorphic to  $M$ .

Analogue of GRR's for groups.

# Structure of $\text{PAut}(\Gamma)$ for graph (digraph, colored graph, multigraph,...) $\Gamma$

**Proposition (R.Jajcay, T.Jajcayova, N.Szakács, M.Szendrei 2021)**

*For any graph  $\Gamma$ , the  $\mathcal{D}$ -classes of  $\text{PAut}(\Gamma)$  correspond to the isomorphism classes of induced subgraphs of  $\Gamma$ , that is, two elements are  $\mathcal{D}$ -related if and only if the subgraphs induced by their respective domains (or images) are isomorphic.*

Partial order for  $\mathcal{D}$ -classes: "subgraph" relation

# Example

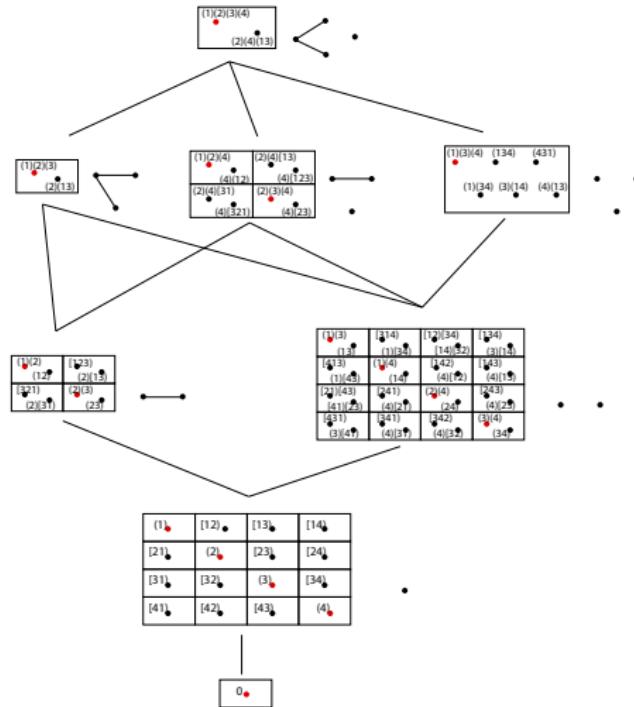


Figure: The Green-class structure of partial graph automorphisms

# When is an inverse monoid of partial permutations the partial automorphism monoid of a graph?

Theorem (R.Jajcay,T.Jajcayova,N.Szakács,M.Szendrei 2021)

Given an inverse submonoid  $S \leq \text{PSym}(X)$ , where  $X$  is a finite set, there exists a graph with vertex set  $X$  whose partial automorphism monoid is  $S$  if and only if the following conditions hold:

1.  $S$  is a full inverse submonoid of  $\text{PSym}(X)$ ,
2. for any compatible subset  $A \subseteq S$  of rank 1 partial permutations, if  $S$  contains the join of any two elements of  $A$ , then  $S$  contains the join of the set  $A$ ,
3. the rank 2 elements of  $S$  form at most two  $\mathcal{D}$ -classes,
4. the  $\mathcal{H}$ -classes of rank 2 elements are nontrivial.

# When is an (abstract) inverse monoid *isomorphic* to the partial automorphism monoid of a graph?

Theorem (R.Jajcay, T.Jajcayova, N.Szakács, M.Szendrei 2021)

*Given a finite inverse monoid  $S$ , there exists a finite graph whose partial automorphism monoid is isomorphic to  $S$  if and only if the following conditions hold:*

1.  $S$  is Boolean,
2.  $S$  is fundamental,
3. for any subset  $A \subseteq S$  of compatible 0-minimal elements, if all 2-element subsets of  $A$  have a join in  $S$ , then the set  $A$  has a join in  $S$ ,
4.  $S$  has at most two  $\mathcal{D}$ -classes of height 2,
5. the  $\mathcal{H}$ -classes of the height 2  $\mathcal{D}$ -classes of  $S$  are nontrivial.

# Symmetry Level of a Graph

## Definition

The **level of symmetry** of a graph  $\Gamma$  of order  $n$  is the ratio between the largest rank of a non-trivial partial automorphism of  $\Gamma$  and its order  $n$ .

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- ▶ A graph  $\Gamma$  admitting a non-trivial automorphism has level of symmetry 1
- ▶ The level of symmetry of  $\Gamma$  is equal to the level of symmetry of its complement

# Symmetry Level of a Graph

**Question 1:** Given  $k \geq 1$ , does there exist a graph  $\Gamma$  of order  $n$  with level of symmetry equal to  $\frac{n-k}{n}$ ?

Computational results (J.Pastorek, V.Cingel 2023+):

- ▶ through exhaustive search: for  $n \leq 10$  all graphs have the level of symmetry at least  $\frac{n-1}{n}$ .
- ▶ he has complete list of graphs on  $n = 11$  vertices with the level of symmetry  $\frac{n-2}{n}$  and there are no graphs with the level of symmetry  $\frac{n-3}{n}$  or less, for  $n = 11$ .
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**Question 2:** What is the minimal level of symmetry of a graph  $\Gamma$  of order  $n$  as a function of  $n$ ?

# Connection to Isomorphism Problem

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- ▶ Result by Rudi Mathon: Determining orbits of  $Aut(\Gamma)$  of a graph is polynomially equivalent to determining  $Aut(\Gamma)$
- ▶ Well-known Weisfeiler-Leman (W-L and k-W-L) Algorithm that approximates orbits  $Aut(\Gamma)$  of a graph. Stable coloring can be computed in time  $O(n^k \log n)$

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- ▶ A partial automorphism  $\varphi$  of  $\Gamma$  is **complete** if there exists no partial automorphism  $\psi$  of  $\Gamma$  that is an extension of  $\varphi$
- ▶ The level of symmetry of  $\Gamma$  of order  $n$  is the ratio between the biggest rank of a *complete* partial automorphism of  $\Gamma$  and  $n$

# Algorithmic aspects of determining $PAut(\Gamma)$

- ▶ A partial automorphism  $\varphi$  of  $\Gamma$  is **complete** if and only if there exists **no** pair of vertices  $u, v$ ,  $u \notin \text{dom}\varphi$  and  $v \notin \text{ran}\varphi$  such that  $u \sim w$  if and only  $v \sim \varphi(w)$ , for all  $w \in \text{dom}\varphi$ .

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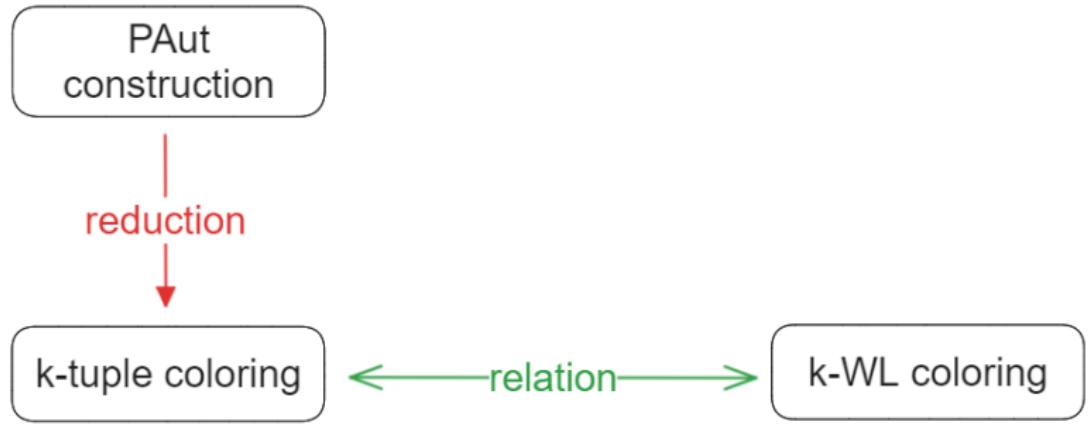
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- ▶ If a partial automorphism  $\varphi$  of  $\Gamma$  is not complete due to a pair  $u, v$ , the partial automorphism  $\tilde{\varphi} : \text{dom}\varphi \cup \{u\} \rightarrow \text{ran}\varphi \cup \{v\}$  is an **extension** of  $\varphi$

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- ▶ Determining whether a partial automorphism is complete and forming all of its extensions of rank increased by 1 is polynomial in the order of  $\Gamma$

# Comparison of the Two Algorithms

- ▶ The algorithm for constructing  $PAut(\Gamma)$  can be adjusted to determine the orbits of  $Aut(\Gamma)$  as follows:  
The vertices  $u, v$  receive different colors at level  $k$  if no partial isomorphism of rank  $k$  can be extended by adding  $u \mapsto v$
- ▶ We are working on comparing the complexity of the two algorithms, and its performance on various graph classes.



# Thank you for listening!

